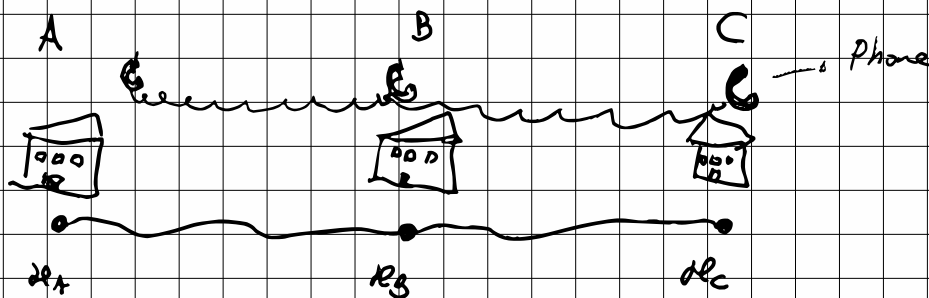


LOCC Local operations and classical communications.

Families of transformations on Composite Systems $\mathcal{H}_A \otimes \mathcal{H}_B \dots$

"Distant-Lab" paradigm



Each party can only influence its own sub-system, but they are allowed to communicate the results of their measurements and other classical information.

Question: given a shared state $|\psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_C$,
what other states can be obtained?

Can we classify entanglement?

2 points entanglement

A & B share $|\psi\rangle = \sum \lambda_i |a_i\rangle |b_i\rangle \in \mathcal{B}(\mathcal{H}_A \otimes \mathcal{H}_B)$

Can they obtain via LOCC Ho state?

$$|\phi\rangle = \sum_i \mu_i |a'_i\rangle |b'_i\rangle$$

Thm (Nielsen)

$|\psi\rangle \succ |\phi\rangle$ iff. $\lambda = (\lambda_1, \dots, \lambda_n)$ **majorizes** $\mu = (\mu_1, \dots, \mu_n)$

Majorization: TFAE

1) λ majorizes μ

2) \exists permutation matrices P_j and prob. distr. $(p_j)_{j=1}^n$ such that

$$\mu = \sum_j p_j P_j \lambda$$

3) \exists doubly stochastic matrix D s.t. $\mu = D\lambda$

$$\mu = D\lambda$$

4) If we reorder

$$\lambda^{\downarrow} = (\lambda_1^{\downarrow}, \dots, \lambda_n^{\downarrow})$$

$$\lambda_1^{\downarrow} \geq \dots \geq \lambda_n^{\downarrow}$$

Then $\sum_{i=1}^q \mu_i^{\downarrow} \leq \sum_{i=1}^q \lambda_i^{\downarrow} \quad \forall q = 1, \dots, n$

OBS 1 Majorization is not stable under tensor products

"Catalytic conversion": $\psi \xrightarrow{\text{Locc}} \phi$ but
 $\psi \otimes \tau \xrightarrow{\text{Locc}} \phi \otimes \tau$

example $\lambda = (0.4, 0.4, 0.1, 0.1)$

$$\mu = (0.5, 0.25, 0.25, 0)$$

$$\tau = (0.6, 0.4)$$

OBS 2 Asymptotics

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2 qubits bipartite state
asymptotically
entanglement entropy

"entanglement gambling"

Start from n copies of $\cos\theta |00\rangle + \sin\theta |11\rangle$

$$\bigotimes_{i=1}^n (\cos\theta |00\rangle + \sin\theta |11\rangle)$$

$$= \sum_{\vec{d}} \cos\theta^{n-|\vec{d}|} \sin\theta^{|\vec{d}|} |\vec{d}, \vec{d}\rangle$$

\vec{d} binary strings
of length n

$$= \sum_{k=0}^n \cos\theta^{n-k} \sin\theta^k \sum_{|\vec{d}|=k} |\vec{d}, \vec{d}\rangle$$

Subspace V_k of dimension $\binom{n}{k}$

Let $P_k = \text{Proj. on Subspace } V_k$

Measuring $\{P_k\}_{k=0}^n$ yields k with prob.

$$P_k = \binom{n}{k} (\cos\theta)^{n-k} (\sin\theta)^k$$

after measuring the new state with a $\binom{n}{k}$ -levels EPR pairs.

expected value is $n \cdot \sin^2\theta$ asymptotically $\left(\frac{\sin^2\theta}{\cos^2\theta}\right)^n \gg \frac{\binom{n}{n/2} \sqrt{2^n}}{\sqrt{2^n}}$

3-partite

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Memu₂



LOCC



GHZ

$$M_{\text{Memu}_2} = \frac{1}{\sqrt{8}} \sum_{i,j,k=0}^1 |i\rangle |j\rangle |k\rangle$$

$M_{\text{Memu}_2} \otimes M_{\text{Memu}_2}$

$$= \frac{1}{8} \sum_{\substack{i,j,k \\ i',j',k'}} |i\rangle |j\rangle |i'\rangle |j'\rangle |k\rangle |k'\rangle$$

$\text{Tr}_A (M_{\text{Memu}_2})$

$$= \frac{1}{8} \sum_{\substack{i,j \\ u,k}} |j\rangle |k\rangle |u\rangle |i\rangle$$

$$= \left(\frac{1}{2} \sum_{i,j} |i\rangle |j\rangle \otimes |i\rangle |j\rangle \right) \otimes \frac{1}{2} \sum_{u,k} |u\rangle |k\rangle \otimes |u\rangle |k\rangle$$

$$= \frac{1}{4} \mathbb{1}_4 \otimes |\text{EPR}\rangle \langle \text{EPR}|$$

$$GHZ_u = \frac{1}{\sqrt{4}} \sum_{i=0}^3 |i\rangle |i\rangle$$

$$|GHZ_u\rangle \langle GHZ_u| = \frac{1}{4} \sum_{\substack{i,j \\ i',j'}} |i\rangle |j\rangle \otimes |i\rangle |j\rangle \otimes |i\rangle |j\rangle$$

$$\text{Tr}_A GHZ_u = \frac{1}{4} \sum_{i=0}^3 |i\rangle |i\rangle \otimes |i\rangle |i\rangle \quad \text{"separable"}$$

THM: ψ, ϕ have same marginal entropies for single sites
are either LU eq. or LOCC-incompatible.

COR 1 LOCC eq. \Leftrightarrow LU eq.

COR 2 if single site marginals ~~have~~ the same entropies but not on larger sites \Rightarrow LOCC incompatible.

Proof $\psi \xrightarrow{\text{local}} \phi \Rightarrow \psi \xrightarrow{\text{LOCC}} \phi$

Step 1 of the protocol maps

$$\psi \rightarrow \psi_i \text{ with prob } p_i$$

$$\rho^{\text{red}} = \sum_i p_i \text{Tr}_A(|\psi_i\rangle\langle\psi_i|) \quad \rho = \sum_i p_i |\psi_i\rangle\langle\psi_i|$$

$$S(\rho_A) = S(\phi_A) \quad \text{and LOCC cannot increase entanglement}$$

$$S(\rho^{\text{red}}) \leq \sum_i p_i S(\text{Tr}_A(|\psi_i\rangle\langle\psi_i|)) \leq S(\rho^{\text{red}}) \leq S(\psi^{\text{red}})$$

$$\Rightarrow \sum_i p_i S(\text{Tr}_A(|\psi_i\rangle\langle\psi_i|)) = S(\sum_i p_i \text{Tr}_A(|\psi_i\rangle\langle\psi_i|))$$

$$\Rightarrow \text{Tr}_A(|\psi_i\rangle\langle\psi_i|) = \text{Tr}_A(|\psi\rangle\langle\psi|)$$

which implies that $|\psi_i\rangle = U_i^\dagger \otimes U_{\text{red}} |\psi\rangle$

3-particles or more

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Stochastic LOCC

"LOCC + post selection" no restriction on success probability

We consider (instead of LOCC) SLOCC (stochastic LOCC):

LOCC but ONLY with a non-zero success probability.

This means that, when one of the parties performs

a measurement $\{P_i\}_{i=1}^k$, we can "post select"

on a given outcome, and "pretend" that it occurs with probability 1.

$|\psi\rangle \xrightarrow{\text{SLOCC}} |\psi\rangle$ iff. \exists local linear operators

L_A, L_B, L_C, \dots such that

$$L_A \otimes L_B \otimes L_C |\psi\rangle = |\psi\rangle \quad (**)$$

if L_A, L_B, L_C are invertible, then $|\psi\rangle \xrightarrow{\text{SLOCC}} |\psi\rangle$, i.e.
the two states are equivalent (**)

(*) is the notion of restriction, which we write
as $|\psi\rangle \triangleright \varphi$.

(**) equivalence classes are orbits of

$$GL_A \otimes GL_B \otimes GL_C.$$

Quantum entanglement SLOCC classification \rightarrow orbits

2-particle: $| \psi \rangle = \sum_{i=1}^n \lambda_i | a_i \rangle | b_i \rangle \geq \sum_{i=1}^n | a_i \rangle | b_i \rangle = \text{SPR}_n$
 $= \langle n, n \rangle$

3-particle

$\triangle_2 = \langle 2 \rangle = \text{GHZ}_2$

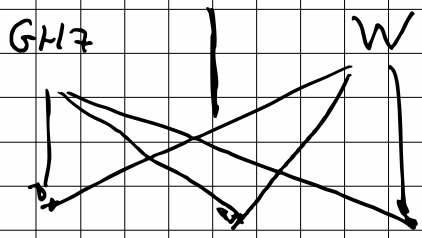
v.s.

$\triangle_W = \frac{1}{\sqrt{3}} (| 001 \rangle + | 010 \rangle + | 100 \rangle)$

are SLOCC - inequivalent.

In fact let $| \psi \rangle$ a 3 qubit state with local ranks r_A, r_B, r_C

multi-partite
entanglement
 $r_A = r_B = r_C = 2$

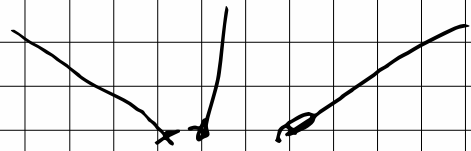


bi-partite
entanglement

$r_A = 2$
 $r_B = r_C = 2$

$r_B = 1$
 $r_A = r_C = 2$

$r_C = 1$
 $r_A = r_B = 2$



product
states

$r_A = r_B = r_C = 1$

$GH_2 \not\cong W$ is equivalent to Segre

$\text{rank } W > 2$

Proof: is direct.

$W \not\cong GH_2$ " " " " Subrank $W > 2$

3-cycle:

$$\hat{c}(W) = 0 \quad \text{and} \quad \hat{c}(GH_2) > 0$$

Assume $|W\rangle = |a_1\rangle|b_1\rangle|c_1\rangle + |a_2\rangle|b_2\rangle|c_2\rangle$

Then $\text{Tr}_A |W\rangle\langle W|$ has range spanned by $|b_1\rangle|c_1\rangle$
and $|b_2\rangle|c_2\rangle$

Let R be a subspace of $\mathbb{C}^2 \otimes \mathbb{C}^2$ which is spanned by

$$V_1 = |v_1^p\rangle \otimes |v_1^q\rangle \quad \text{and} \quad V_2 = |v_2^p\rangle \otimes |v_2^q\rangle$$

Then either R contains all product vectors,

Or R contains no other product vectors.

Case 1 exclude since local ranks are more 2.

$$R = \text{Span} \{ |01\rangle + |10\rangle, |00\rangle \} = \left\{ \begin{pmatrix} \alpha \\ \beta \\ 0 \\ 0 \end{pmatrix} \mid \alpha, \beta \in \mathbb{C} \right\}$$

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3-tangle is $\left| \begin{array}{c} \text{Codes} \\ \text{Hyper-determinant} \end{array} \right|$

dim orbit GHZ = 7

dim orbit ψ = 6

dim orbit biseparable = 4

dim completely separable = 3

maximal dimension of 3-qubits

